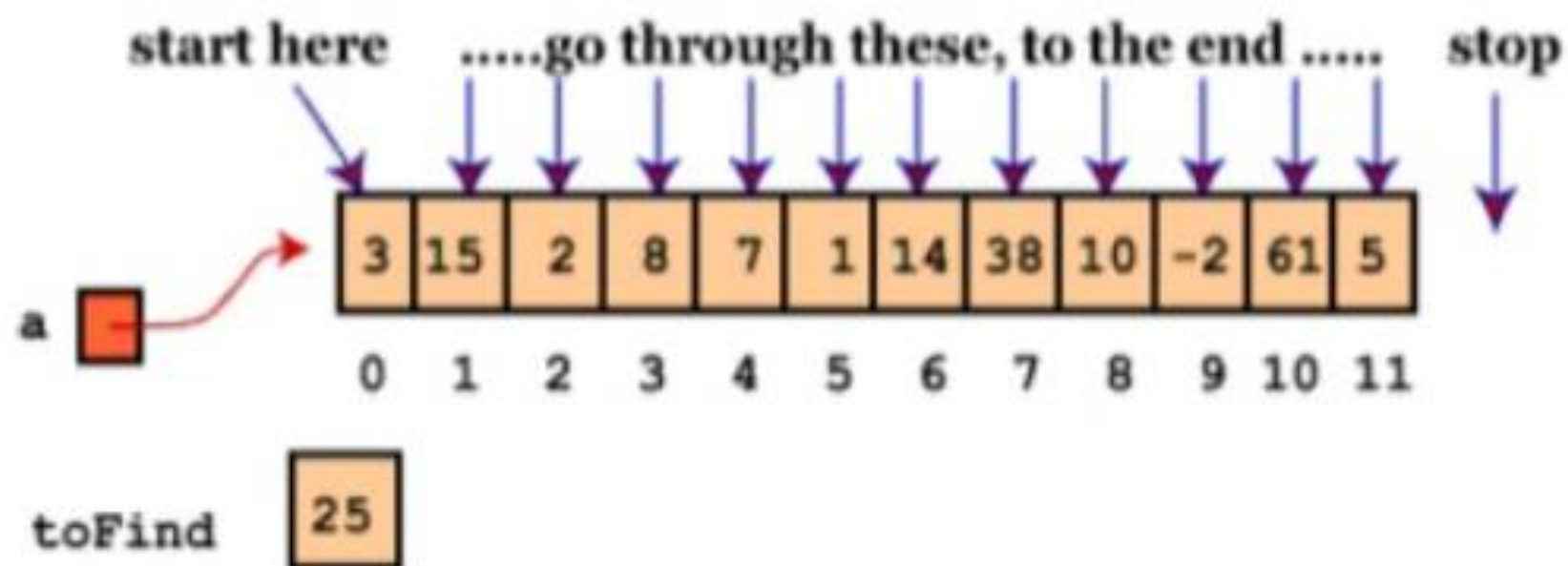


Linear Search
&
Binary Search



Every item is checked but no match is found till the end of the data collection

- Find 37?

0	1	2	3	4	5	6	7	8
20	35	37	40	45	50	51	55	67
↑	↑	↑						
≠	≠	=						
		Return 2						

Found a match at index 2

- procedure linear_search (list, value)

```
for each item in the list
  if match item == value
    return the item's location
  end if
end for
```

```
end procedure
```

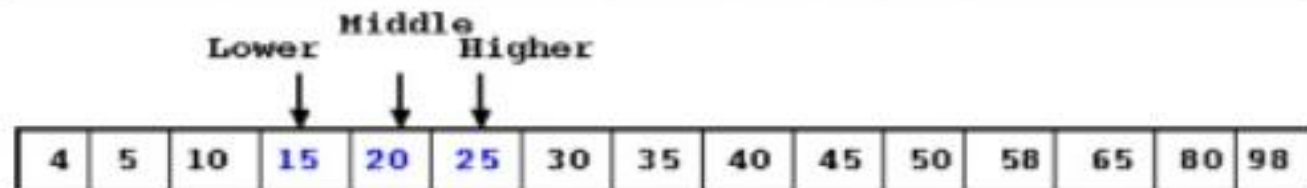
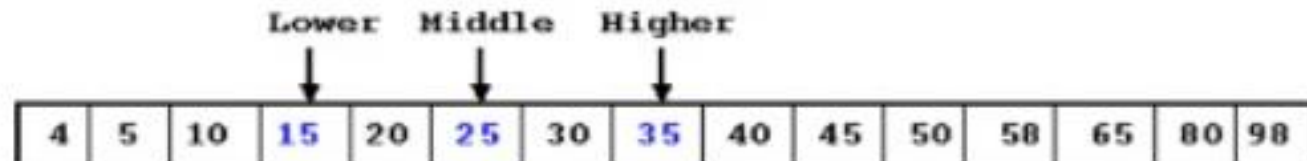
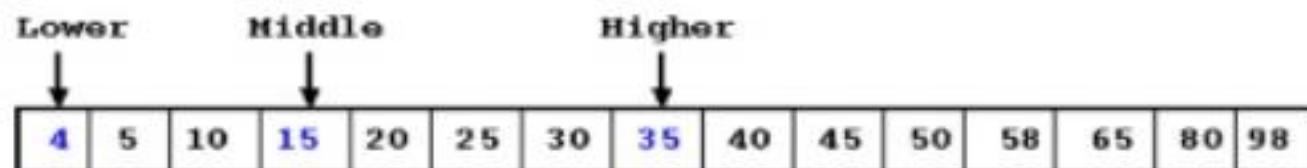
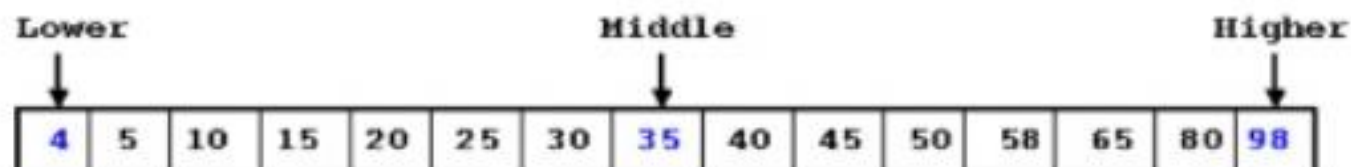
Adv. & Disadv. Of LS

- Advantages
 - Easiest to understand and implement
 - No sorting required
 - Suitable for small list sizes
 - Works fine for small number of elements
- Disadvantages
 - Time inefficient as compared to other algorithms
 - Not suitable for large-sized lists
 - Search time increases with number of elements

Binary Search (BS)

- Binary Search is a Divide and Conquer algorithm
- Binary search algorithm finds the position of a target value within a sorted array
- A more efficient approach than Linear Search because Binary Search basically reduces the search space to half at each step

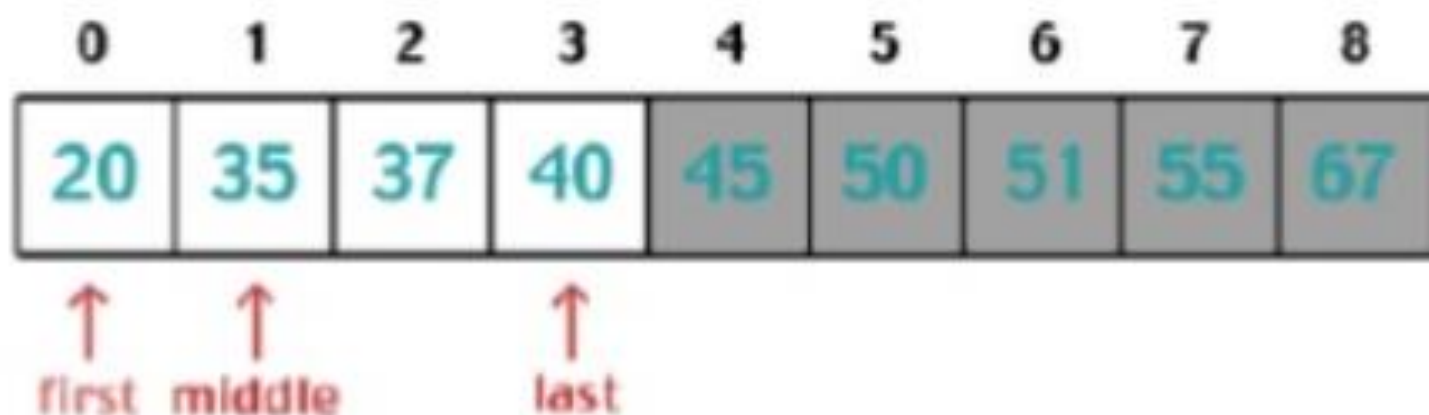
Graphical Illustration of BS



- Find 37?
 1. Sort Array.

0	1	2	3	4	5	6	7	8
20	35	37	40	45	50	51	55	67

Repeat 2. Calculate $\text{middle} = (\text{low} + \text{high}) / 2$.
 $= (0 + 3) / 2 = 1$.

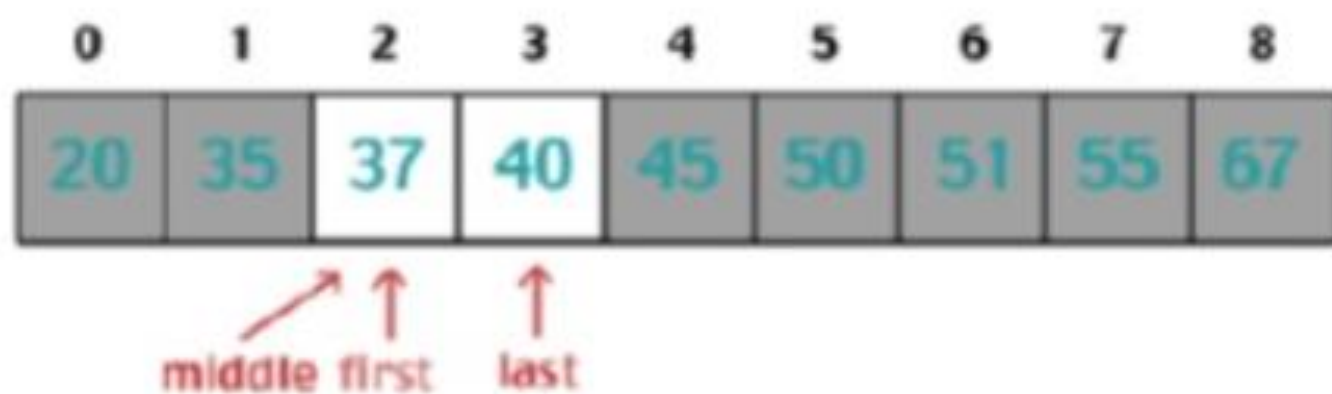


If $37 == \text{array}[\text{middle}] \rightarrow \text{return middle}$

Else if $37 < \text{array}[\text{middle}] \rightarrow \text{high} = \text{middle} - 1$

Else if $37 > \text{array}[\text{middle}] \rightarrow \text{low} = \text{middle} + 1$

Repeat 2. Calculate $\text{middle} = (\text{low} + \text{high}) / 2$.
 $= (2 + 3) / 2 = 2$.



If $37 == \text{array}[\text{middle}] \rightarrow$ return middle

Else if $37 < \text{array}[\text{middle}] \rightarrow$ high = middle - 1

Else if $37 > \text{array}[\text{middle}] \rightarrow$ low = middle + 1

Binary Search

- With each test that fails to find a match, the search is continued with one or other of the two sub-intervals, each at most half the size
- If the original number of items is N then after the first iteration there will be at most $N/2$ items remaining, then at most $N/4$ items, and so on
- In the worst case, when the value is not in the list, the algorithm must continue iterating until the list is empty

Pseudocode

```
Procedure binary_search
  A ← sorted array
  n ← size of array
  x ← value to be searched

  Set lowerBound = 1
  Set upperBound = n

  while x not found
  if upperBound < lowerBound
    EXIT: x does not
    exists.
```

```
    set midPoint =
    lowerBound + ( upperBound - lowerBound )/2

    if A[midPoint] < x
      set lowerBound = midPoint + 1

    if A[midPoint] > x
      set upperBound = midPoint - 1

    if A[midPoint] = x
      EXIT: x found at location midPoint
    end while

  end procedure
```

Iterative binary search

```
int begin = 0;  
int last = array.Length - 1;  
int mid = 0;
```

Part #1 Initialize pointers

```
while (begin <= last) {  
    mid = (begin + last) / 2;  
    if (array[mid] < x) {  
        begin = mid + 1;  
    }  
    else if (array[mid] > x) {  
        last = mid - 1;  
    }  
    else {  
        return mid;  
    }  
}  
  
return -1;
```

Part #2 Search